

Propositional Calculus - Truth Table

Logic Gate Terminology		NOR		inversion		inversion		~XOR		NAND		equivalence		implication		implication		tautology	
p	q	contradiction	non-implication		non-implication		non-equivalence		AND		XOR		amplifier		amplifier		OR		
1	1	F	F	F	F	F	F	F	F	T	T	T	T	T	T	T	T	T	
1	0	F	F	F	F	T	T	T	T	F	F	F	F	T	T	T	T	T	
0	1	F	F	T	T	F	F	T	T	F	F	T	T	F	F	T	T	T	
0	0	F	T	F	T	F	T	F	T	F	T	F	T	F	T	F	T	T	
OR function		~(p v ~p)	~(p v ~q)		~(~p v q)		~(p v ~q) v ~(~p v q)		~(~p v ~q)		q v q		p v p		~(p v ~q)		~(p v ~p)		
		~(p v q)		~p		~q		~(p v ~q)		~(~p v ~q) v ~(p v q)		~(p v q)		~(p v ~q)		~(p v ~p)			
Conditional function (1)		~(~p ⊃ ~p)	~(~p ⊃ ~q)		~(p ⊃ q)		~(p ⊃ q) ⊃ ~(~p ⊃ ~q)		~(p ⊃ ~q)		~(q ⊃ q)		~(~p ⊃ p)		~(p ⊃ ~q)		~(p ⊃ ~p)		
Implication rule		~(~p ⊃ q)	~(p ⊃ ~p)		~(q ⊃ ~q)		~(p ⊃ ~q)		~(p ⊃ ~q) ⊃ ~(~p ⊃ p)		~(p ⊃ q)		~(p ⊃ ~q)		~(p ⊃ ~q)		~(p ⊃ ~p)		
Conditional function (2)		~(p ⊃ p)	~(q ⊃ p)		~(~q ⊃ ~p)		~(p ⊃ ~q) ⊃ ~(p ⊃ q)		~(q ⊃ ~p)		(same)		(same)		~(q ⊃ p)		~(q ⊃ p)		
Transposition rule		~(~q ⊃ p)	(same)		(same)		~(q ⊃ ~p)		~(p ⊃ q) ⊃ ~(p ⊃ ~q)		~(q ⊃ ~p)		~(p ⊃ q)		~(q ⊃ p)		~(p ⊃ p)		
AND function		~(p & p)	~(p & q)		~(p & ~q)		~(p & q) & ~(~p & ~q)		~(p & q)		~(q & q)		~(p & p)		~(p & ~q)		~(~p & ~q)		
(de Morgan's rule)		~(p & ~q)	~(p & p)		~(q & q)		~(p & q)		~(p & ~q) & ~(~p & ~q)		~(p & ~q)		~(p & ~q)		~(p & ~q)		~(~p & ~p)		
Equivalence function		p ≡ ~p ⊃ ~p	~p ≡ ~p ⊃ q		p ≡ ~p ⊃ ~q		~p ≡ ~q		p ≡ ~p ⊃ q		q ≡ ~q ⊃ q		p ≡ ~p ⊃ p		~(~p ≡ ~p ⊃ ~q)		~(~p ≡ ~p ⊃ ~q)		
(exclusive OR)		~p ≡ ~p ⊃ ~q	~(p ≡ ~p ⊃ p)		~(q ≡ ~q ⊃ q)		~(p ≡ ~p ⊃ q)		~(p ≡ q)		~(p ≡ ~p ⊃ ~q)		~(p ≡ ~p ⊃ q)		~(p ≡ ~p ⊃ q)		~(p ≡ ~p ⊃ q)		
Equivalence function		~p ≡ ~p ⊃ ~q		~(p ≡ ~p ⊃ p)		~(q ≡ ~q ⊃ q)		~(p ≡ ~p ⊃ q)		~(p ≡ q)		~(p ≡ ~p ⊃ ~q)		~(p ≡ ~p ⊃ q)		~(p ≡ ~p ⊃ q)		~(p ≡ ~p ⊃ q)	
NAND function		~(~p p)	~(~p q)		~(p ~q)		~(~p ~q)		~(p q)		~(q q)		~(p p)		~(p ~q)		~(p ~q)		
(Sheffer's rule)		~(~p ~q)	~(p p)		~(q q)		~(p ~q)		~(p q)		~(p q ~p ~q)		~(p q)		~(p q ~p ~q)		~(p q)		
NOR function		~(p ↓ ~p)	~(p ↓ ~q)		~(p ↓ q)		~(p ↓ q) ↓ ~p ↓ ~q		~(p ↓ ~q)		~(q ↓ q)		~(p ↓ ~p)		~(p ↓ ~q)		~(p ↓ q)		
(Joint denial)		~(p ↓ q)	~(p ↓ p)		~(q ↓ q)		~(~p ↓ ~q)		~(~p ↓ ~q)		~(~p ↓ ~q)		~(~p ↓ ~q)		~(~p ↓ ~q)		~(~p ↓ ~q)		

Symbols of Propositional Calculus

Symbol	Term	1st. Meaning	2nd. Meaning	Logic Gate
~	negation	not		not
&	conjunction	and		and
v	disjunction	or	(from Latin vel)	(inclusive) or
l	incompatible	at least one is false	cannot both be true	nand
↓	joint denial	both false	neither is true	nor
≡	biconditional	if and only if	both true or both false	exclusive or
⊃	conditional	→ if-then	p implies q	implication

Axioms

$\neg(p \vee p) \supset p$	1st axiom of propositional calculus
$p \supset \neg(p \vee \neg p)$	2nd axiom of propositional calculus
$\neg(p \vee q) \supset \neg(q \vee p)$	3rd axiom of propositional calculus
$\neg(p \supset q) \supset \neg(\neg r \vee p) \supset \neg(r \vee q)$	4th axiom of propositional calculus
$\neg(p \vee \neg p)$	<i>Tertium-non-datur</i> - the excluded middle even if p is indeterminate .

Rules of Propositional Calculus

Joining Rule	If p and q are theorems then $\neg(p \& q)$ is a theorem.
Separation Rule	If $\neg(p \& q)$ is a theorem then both p and q are theorems
Double tilde rule	$\neg\neg$ can be deleted or inserted into any theorem (providing it is well-formed)
Deduction Rule	If q can be derived when p is <u>assumed</u> then $\neg(p \supset q)$ is a theorem (a fantasy).
Detachment Rule	If p and $\neg(p \supset q)$ are both theorems, then q is a theorem - the <i>modus ponens</i>
(or Derivation Rule)	<i>Deduction and Detachment are complementary</i>
Carry-over Rule	Inside a fantasy, any theorem from the reality one level higher can be used.
Implication Rule	$\neg(p \vee q)$ and $\neg(\neg p \supset \neg q)$ are interchangeable.(Hofstadter terms the Switcheroo)
Transposition Rule	$\neg(p \supset q)$ and $\neg(\neg q \supset \neg p)$ are interchangeable.(Hofstadter terms Contrapositive)
de Morgan's Rule	$\neg(p \& \neg q)$ and $\neg(\neg p \vee q)$ are interchangeable

Proofs

$\neg(\neg p \supset q) \& \neg(q \supset r) \supset \neg(p \supset r)$	Proof by hypothetical syllogism (if p implies q and q implies r then p implies r)
$\neg(\neg p \vee q) \& \neg q \supset p$	Proof by disjunctive syllogism (If either p or q true but q false then p must be true)
$\neg(p \vee q) \& \neg p \supset \neg q$	<i>This is an expanded version of proof by disjunctive syllogism</i>
$\neg(p \& \neg p) \supset q$	Proof by indirect reasoning <i>Modus ponens</i> (if p then q but p is true so q must be true)
$\neg(\neg p \supset q) \& \neg q \supset \neg p$	Proof by detachment <i>Modus tollens</i> (if p then q but if q is false then so must p)
$\neg(\neg p \supset q) \& \neg(\neg p \supset \neg q) \supset p$	Proof by contradiction <i>Reductio ad absurdum</i> . (If $\neg p$ leads to a contradiction, p must be true)
$\neg(p \vee q) \& \neg(\neg p \supset r) \& \neg(q \supset r) \supset r$	Proof by cases. (if either p or q is true and either imply r, then r must also be true)

References

Forever Undecided - a puzzle guide to Godel : Raymond Smullyan : <i>Oxford Press</i>
Mathematical Logic : R.L. Goodstein : <i>Leicester University Press</i>
Godel, Escher, Bach - an eternal golden braid : Douglas R. Hofstadter : <i>Penguin Books</i>
University of Bath Sequential Systems Tutor Notes from 1972

Equivalences

Tautologies	$p = p \vee p$	$p = p \vee p$
Commutative	$p \& q = q \& p$	$p \vee q = q \vee p$
Associative	$p \& (q \& r) = (p \& q) \& r$	$p \vee (q \vee r) = (p \vee q) \vee r$
Distributive	$p \& (q \vee r) = p \& q \vee p \& r$	$p \vee (q \& r) = p \vee q \& p \vee r$
q v ~q can be deleted from any expansion		